THINK-TAC-TOE: WHEN ARE PUZZLES SOLVABLE?



Susa Stonedahl & Forrest Stonedahl

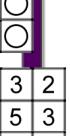
MathFest: August 5, 2011

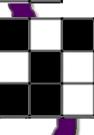
Motivation

Co-taught a middle school math enrichment program

- Students like puzzles
- Experienced and inexperienced mathematicians are on more even ground when facing a new puzzle



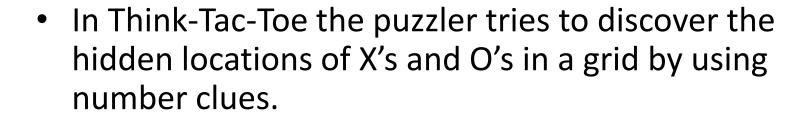




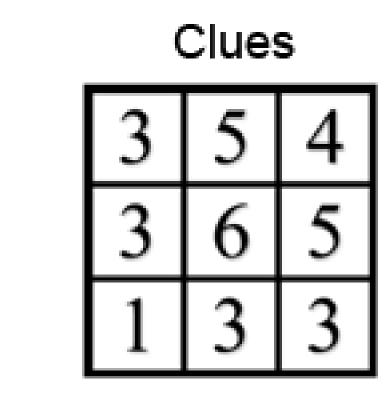


2 3 2 4 5 3 3 3 2

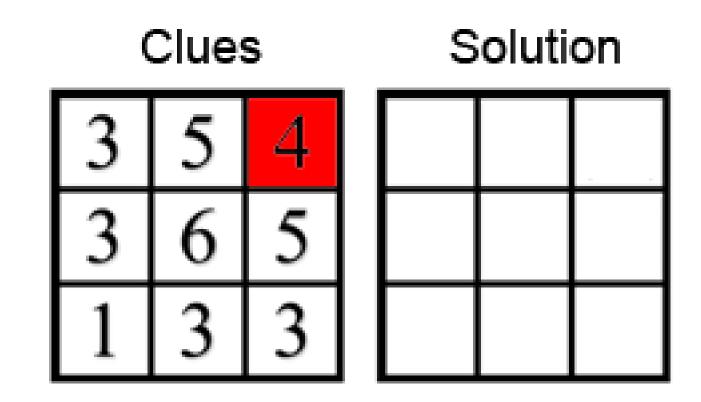
Rules of Think-Tac-Toe



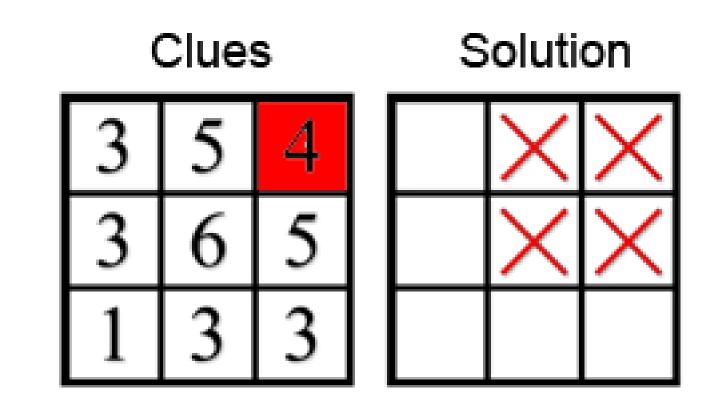
- The number in each square tells you the number of X's in that square's neighborhood.
 - A square's neighborhood is made up of the square itself and any squares that it shares an edge or a corner with.



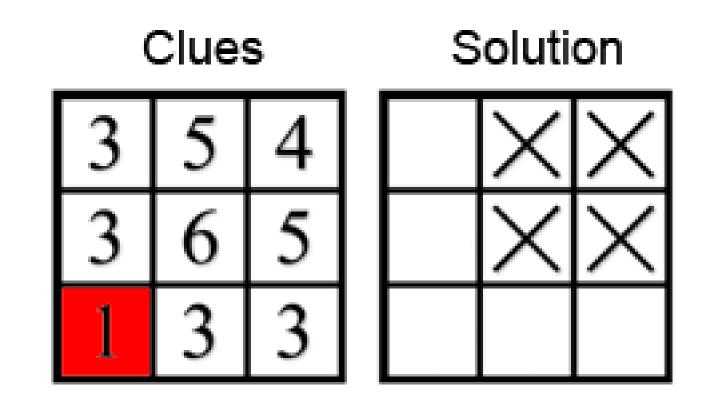
Solution



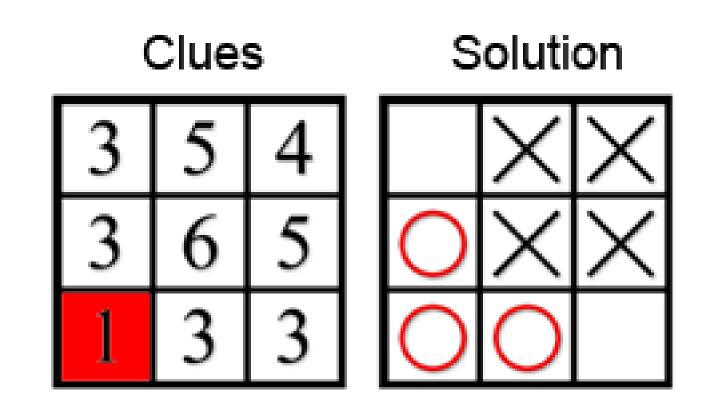
If we look at the 4....



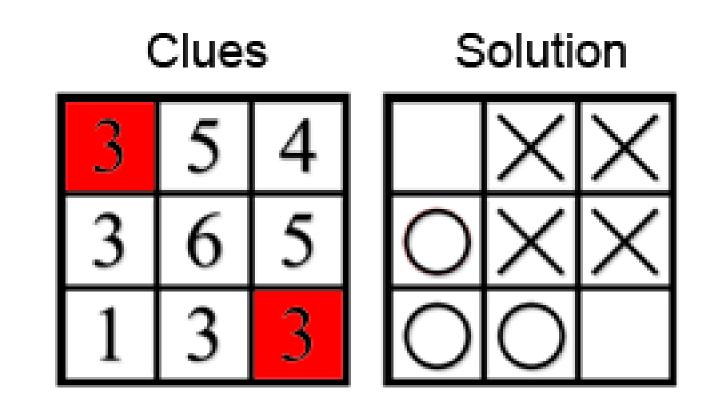
The whole neighborhood has X's



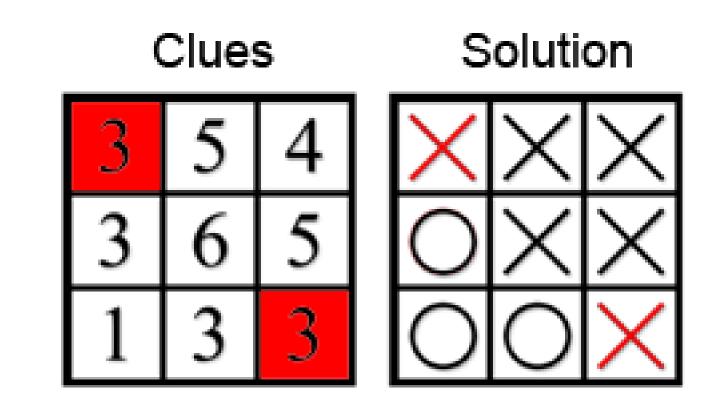
If we look at the 1...



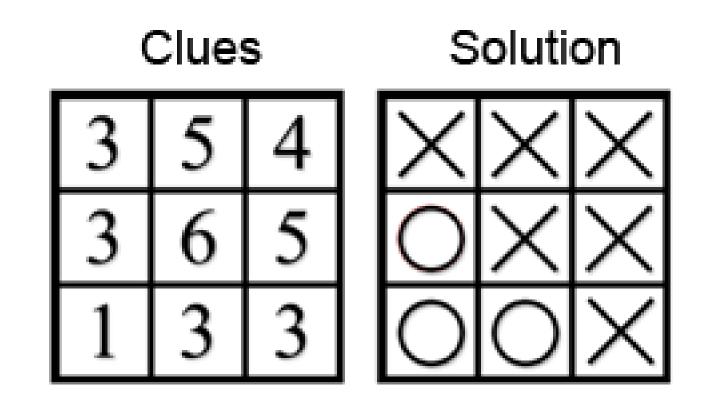
The neighborhood already has it's "1"



If we look at these 3's...

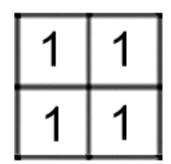


They each need another X.



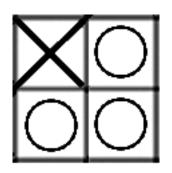
The puzzle is solved!

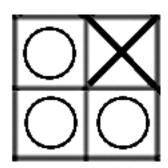
But NOT All Puzzles Are Solvable...

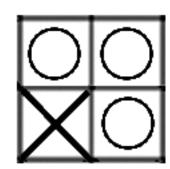


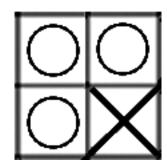
This puzzle could have originated as any of these four solutions, so it's not solvable.

(Puzzle creation operation is not invertible!)



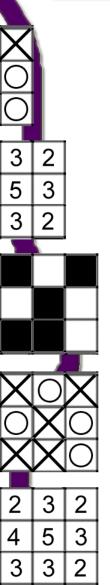




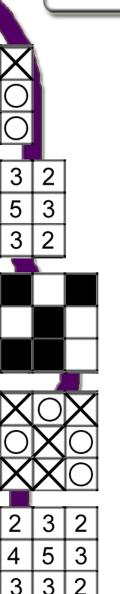


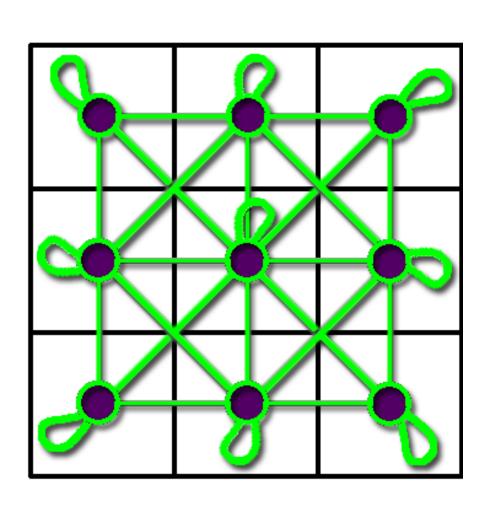
Our Question is...

Which grid sizes are always solvable?



We can treat the puzzle grid as a graph





3×3 Adjacency Matrix...

T _(1,1)	T _(1,2)	T _(1,3)
T _(2,1)	T _(2,2)	T _(2,3)
T _(3,1)	T _(3,2)	T _(3,3)

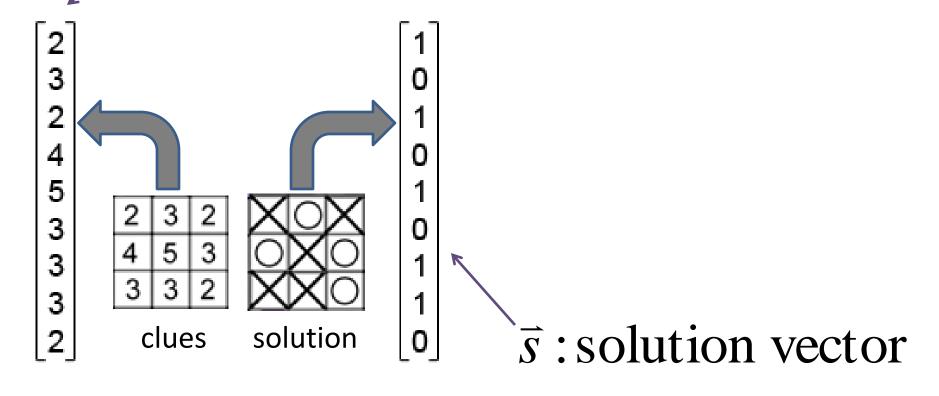
 $T_{(1,1)}T_{(1,2)} T_{(1,3)} T_{(2,1)}T_{(2,2)} T_{(2,3)} T_{(3,1)} T_{(3,2)} T_{(3,3)}$ $T_{(1,1)}$ T_(1,2) 1 1 1 1 1 1 0 $T_{(1,3)}$ $T_{(2,1)}$ 1 1 0 1 1 0 1 T_(2,2) 1 1 1 1 1 1 1 1 T_(2,3) 0 1 1 0 1 1 0 0 0 0 1 1 0 1 $T_{(3,1)}$ $T_{(3,2)}$

T: locations in the grid

A: corresponding adjacency matrix

3×3 Solution and Clue Vectors...

 \vec{c} : clues vector (#'s from 0 to 9)



(0's for O's and 1's for X's)

Solvability

creating puzzle...

$$A\vec{s} = \vec{c}$$

A: adjacency matrix

 \vec{s} : solution vector

 \vec{c} : clues vector

Solvability

creating puzzle...

$$A\vec{s} = \vec{c}$$

solving puzzle...

$$\vec{s} = A^{-1}\vec{c}$$

A: adjacency matrix

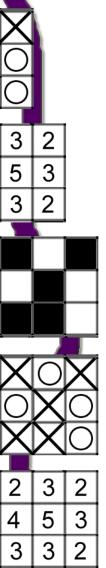
 \vec{s} : solution vector

 \vec{c} : clues vector

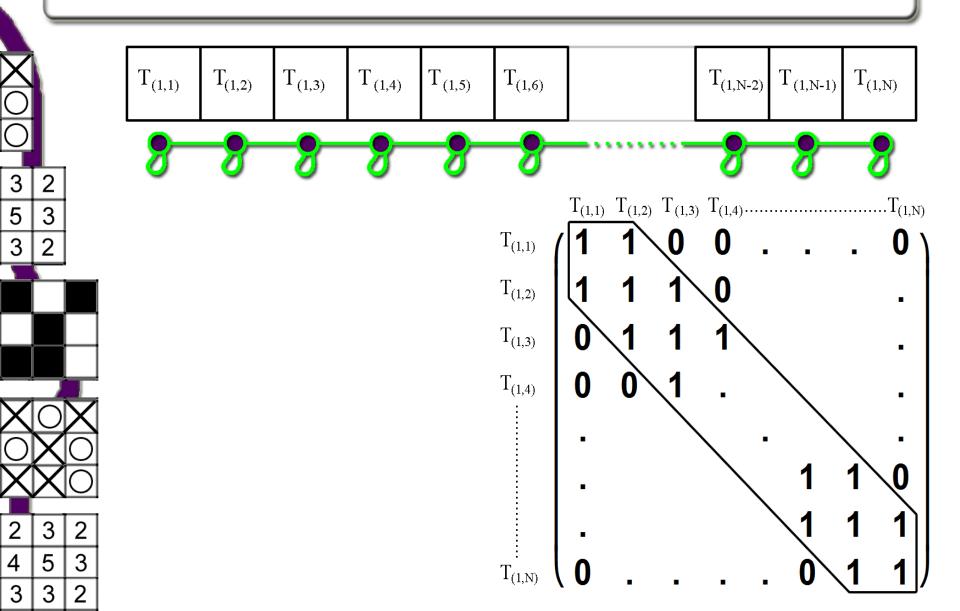
If the corresponding adjacency matrix, A, is invertible, then the puzzle is solvable!

When is the 1×N Puzzle Solvable?

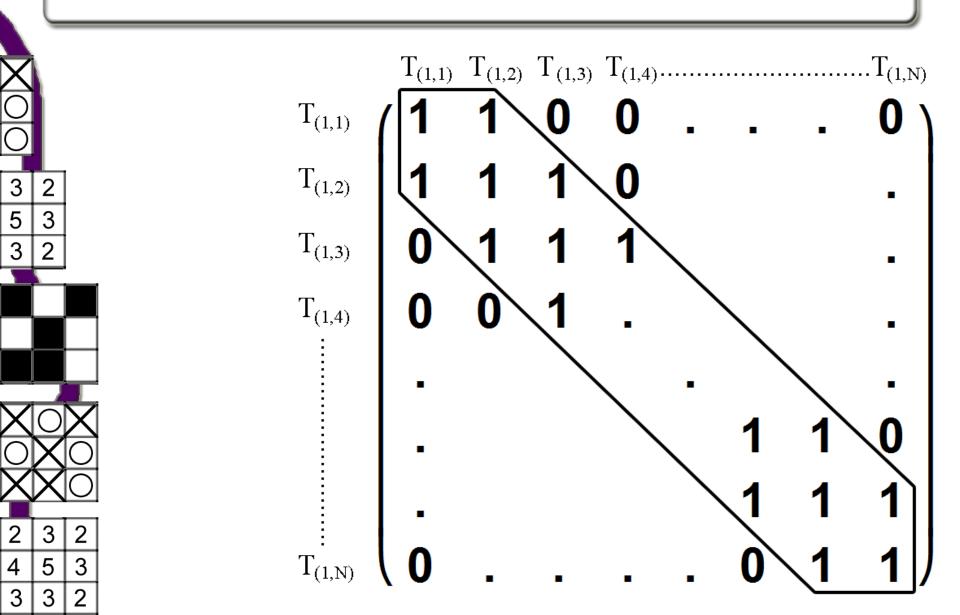
 Our goal is to discover when an M×N matrix puzzle is solvable, but let's solve a simpler problem first.



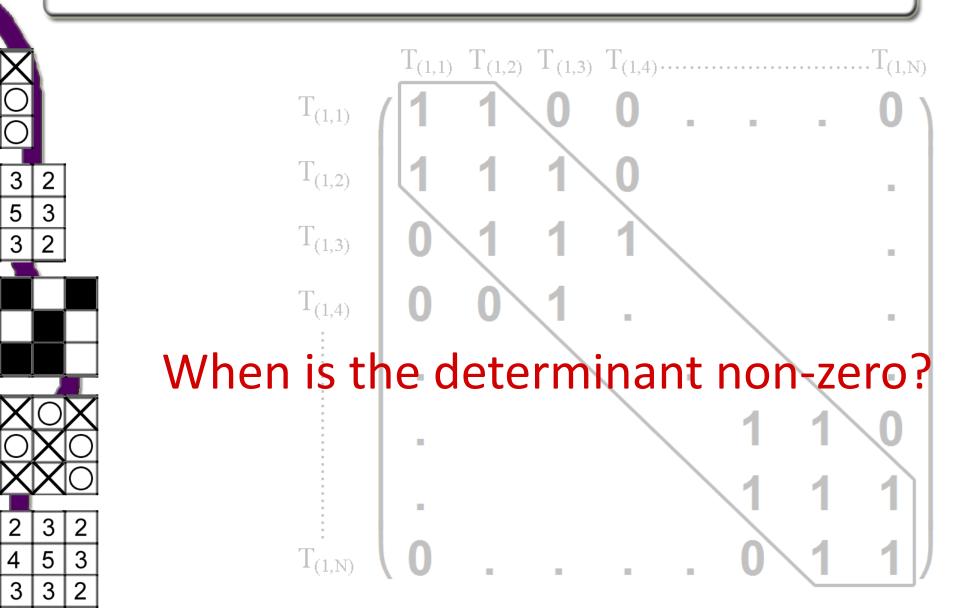
When is the 1×N Puzzle Solvable?



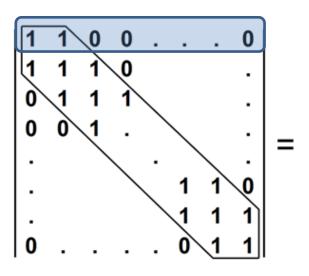
When is the 1×N Adjacency Matrix Invertible?



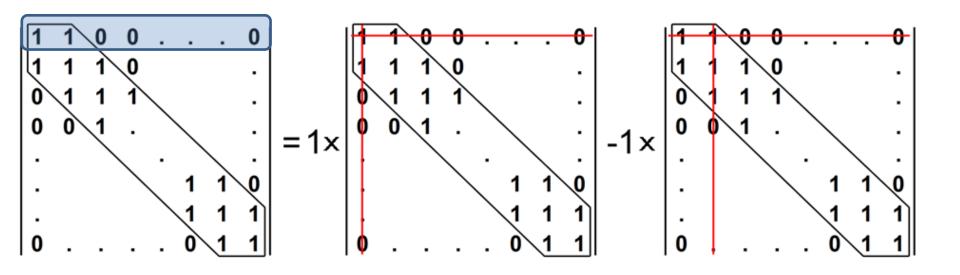
When is the 1×N Adjacency Matrix Invertible?



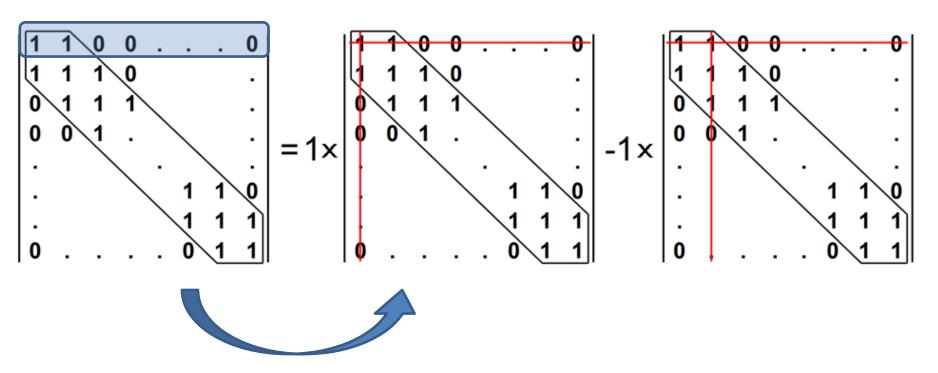
Determinant of a 1×k adjacency matrix..



Determinant of a 1×k adjacency matrix..

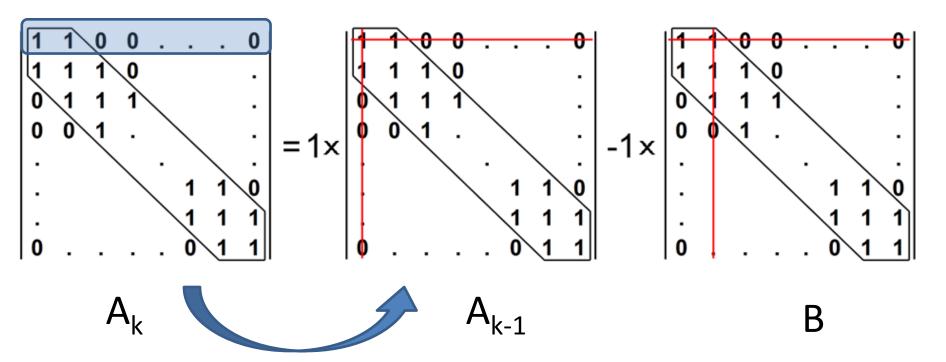


Determinant of a 1×k adjacency matrix...



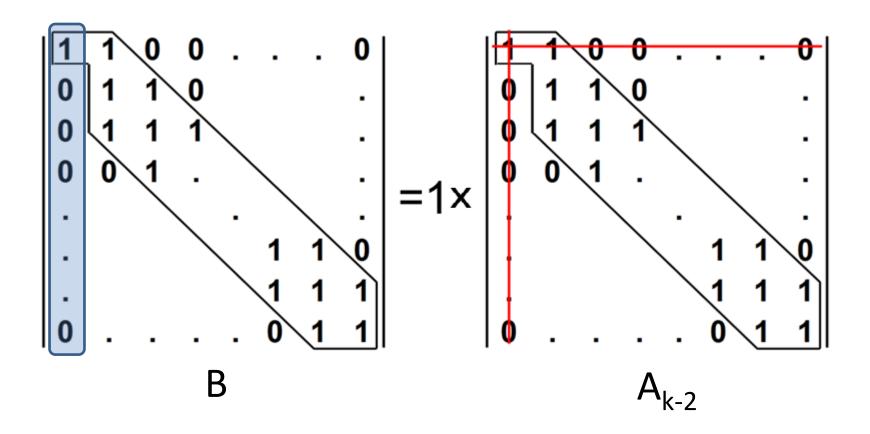
The same form of a matrix, but with a different number of rows and columns

Determinant of a 1×k adjacency matrix...



The same form of a matrix, but with a different number of rows and columns

Expanding again...



Expanding by the first column.

Determinant of a 1×k adjacency matrix...

$$\begin{vmatrix} 1 & 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 1 & 0 & & \dots \\ 0 & 1 & 1 & 1 & & \dots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 1 & & \dots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0$$

$$det(A_k)=det(A_{k-1})-det(A_{k-2})$$

 $det(A_k) = det(A_{k-1}) - det(A_{k-2})$





X	\bigcirc	X
\bigcirc	X	O
X	X	

2	3	2
4	5	3

3 3 2

$$det(A_k)=det(A_{k-1})-det(A_{k-2})$$

$$det(A_{k-1}) = det(A_{k-2}) - det(A_{k-3})$$



- 2 3 24 5 3
- 3 3 2

$$det(A_k) = det(A_{k-1}) - det(A_{k-2})$$

$$det(A_{k-1}) = det(A_{k-2}) - det(A_{k-3})$$

$$det(A_k) = det(A_{k-2}) - det(A_{k-3}) - det(A_{k-2})$$





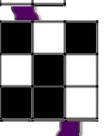


- 2 3 2 4 5 3
- 3 3 2

$$det(A_k)=det(A_{k-1})-det(A_{k-2})$$

$$det(A_{k-1}) = det(A_{k-2}) - det(A_{k-3})$$

$$det(A_k) = det(A_{k-2}) - det(A_{k-3}) - det(A_{k-2})$$



X	\bigcirc	X
O	X	O
X	X	$oxed{\Box}$

 2
 3
 2

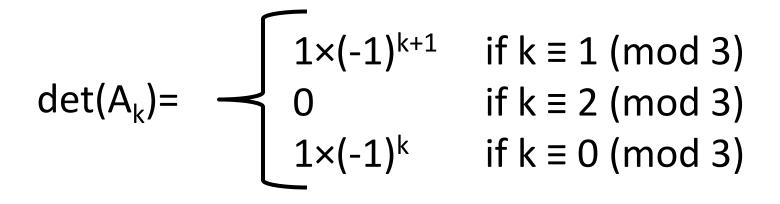
 4
 5
 3

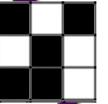
$$det(A_k)=det(A_{k-1})-det(A_{k-2})$$

$$det(A_{k-1}) = det(A_{k-2}) - det(A_{k-3})$$

$$det(A_k) = -det(A_{k-3})$$

Proof by Strong Induction...





Base cases:

$$det(A_1)=1$$

$$det(A_2)=0$$

$$det(A_3)=-1$$

Given:

$$det(A_k) = -det(A_{k-3})$$

Lemma

 $1\times N$ puzzles are solvable iff $N \neq 2 \pmod{3}$





2	3	2
4	5	3
$\overline{}$	$\overline{}$	

When is M×N Solvable?



Now that we've solved the 1×N case, let's solve the more general M×N case!

Is the 3×3 Solvable?

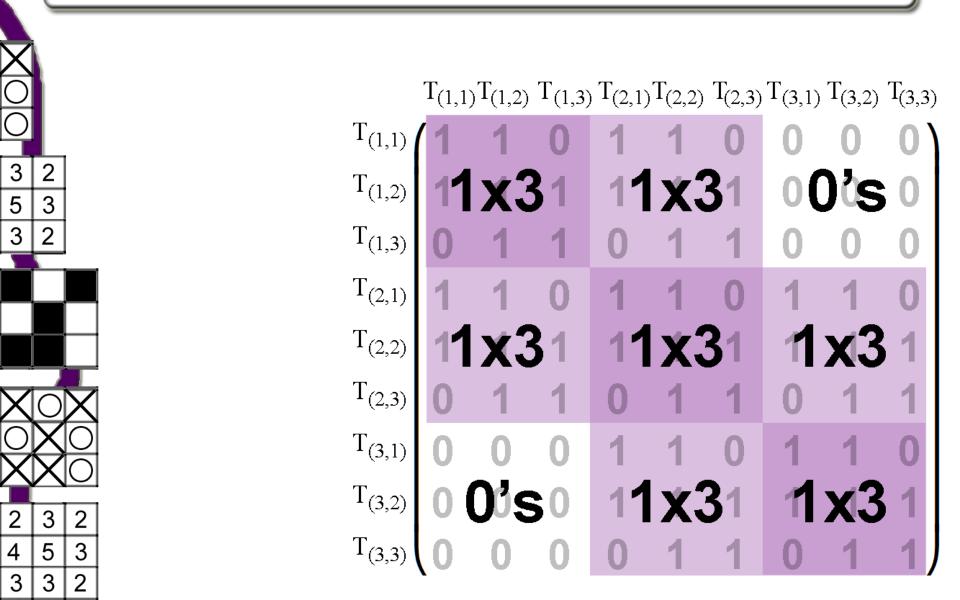
Yes, row reducing this matrix yields the identity

	$T_{(1,1)}$	$1_{(1,2)}$	$T_{(1,3)}$	$T_{(2,1)}$	$T_{(2,2)}$	$T_{(2,3)}$	$T_{(3,1)}$	$T_{(3,2)}$	$T_{(3,3)}$
$T_{(1,1)}$	(1	1	0	1	1	0	0	0	0
$T_{(1,1)}$ $T_{(1,2)}$	1	1	1	1	1	1	0	0	0
$T_{(1,3)}$	0	1	1	0	1	1	0	0	0
$T_{(2,1)}$	1	1	0	1	1	0	1	1	0
$T_{(2,2)}$	1	1	1	1	1	1	1	1	1
$T_{(2,3)}$	0	1	1	0	1	1	0	1	1
$T_{(3,1)}$	0	0	0	1	1	0	1	1	0
$T_{(3,2)}$	0	0	0	1	1	1	1	1	1
$T_{(3,3)}$	0)	0	0	0	1	1	0	1	1

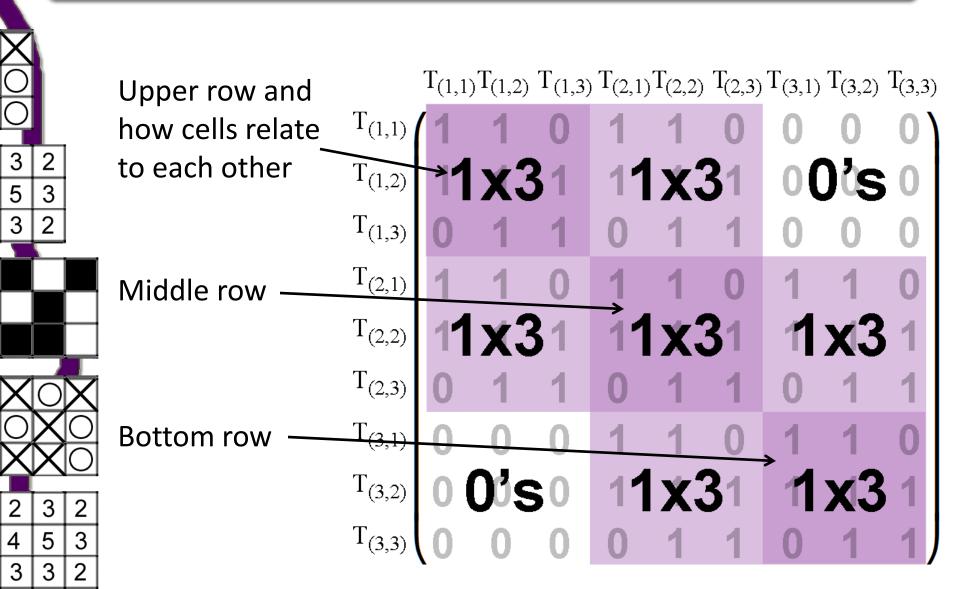
Looking for patterns...

		$^{1}(1,1)$	$)^{1}(1,2)$	$^{1}(1,3)$	1(2,1]	$)^{1}(2,2)$	$I_{(2,3)}$	$^{1}(3,1)$	1(3,2)	$^{1}(3,$
The same	$T_{(1,1)}$	(1	1	0	1	1	0	0	0	0
3×3 block is	$T_{(1,2)}$	1	1	1	1	1	1	0	0	0
repeated	$T_{(1,3)}$	0	1	1	0	1	1	0	0	0
	$T_{(2,1)}$	1	1	0	1	1	0	1	1	0
	$T_{(2,2)}$	1	1	1	1	1	1	1	1	1
	$T_{(2,3)}$	0	1	1	0	1	1	0	1	1
	$T_{(3,1)}$	0	0	0	1	1	0	1	1	0
	$T_{(3,2)}$	0	0	0	1	1	1	1	1	1
	$T_{(3,3)}$	n	0	O	0	1	1	n	1	1

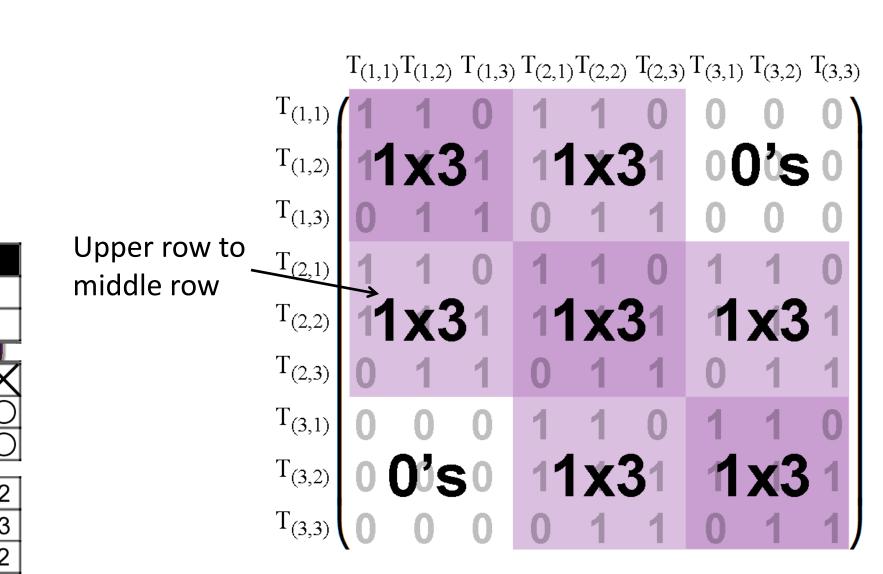
The blocks are in fact the 1×3 A matrix



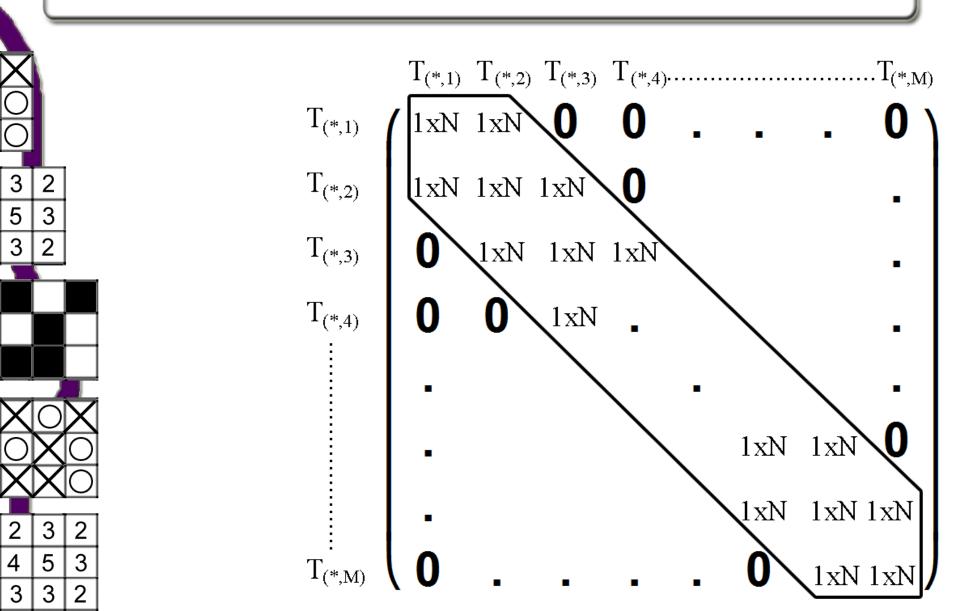
The blocks are in fact the 1×3 A matrix



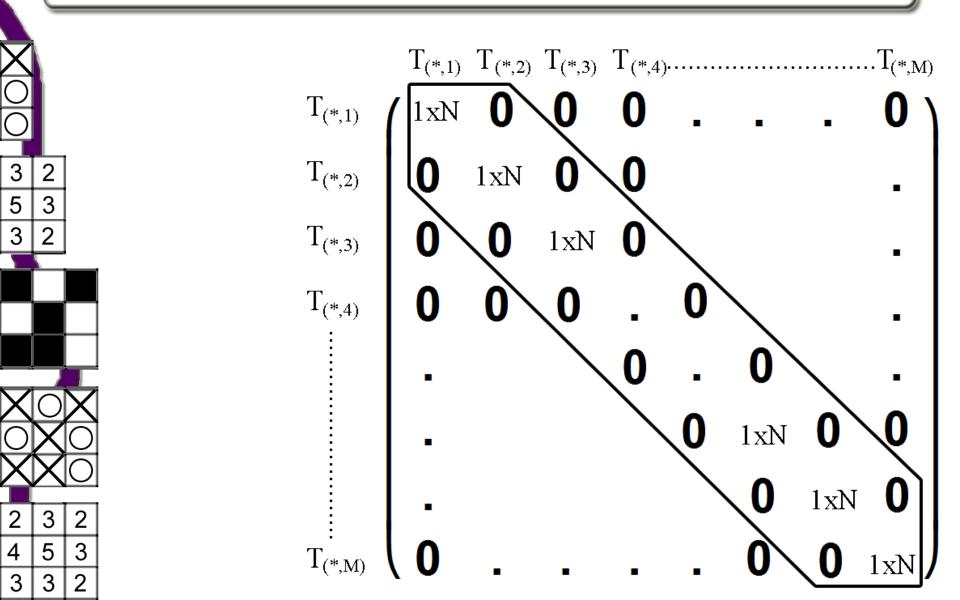
The blocks are in fact the 1×3 A matrix



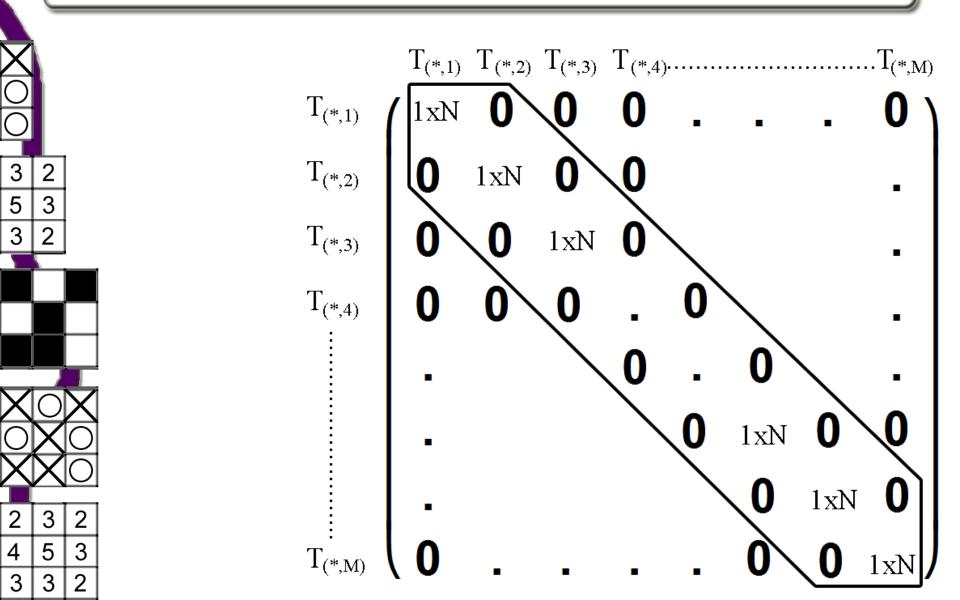
And the 1×N blocks for an M×N...



The M×N can be row reduced like this if the 1×M can be row reduced... [M ≠2 (mod 3)]



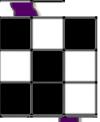
Each 1×N can be row reduced if N ≠2 (mod 3)



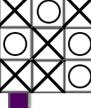
End of the Proof



So any M×N puzzle is solvable if the 1×N and 1×M versions are solvable.



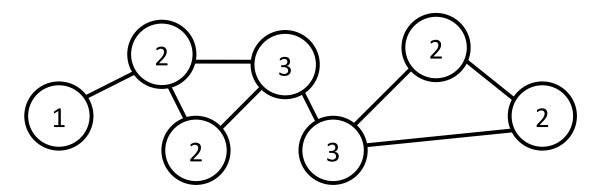
By our lemma, this is true whenever N ≠2 (mod 3) and M ≠2 (mod 3)

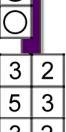


QED

Open Problems / Future Work

- Alternative lattice structures?
- Alternative neighborhoods:
 - e.g. don't count yourself
- Can we characterize when a puzzle is uniquely solvable when played on an arbitrary graph?











Take Think-Tac-Toe with you!

- Give the puzzles a try they're fun!
- Give them to students
- Solve an open question from the previous slide
- Develop a new variation
- Something completely different

And tell us about it – we'd love to hear from you!

http://www.stonedahl.com/thinktactoe/

QUESTIONS?

Extra slides follow



3 25 3

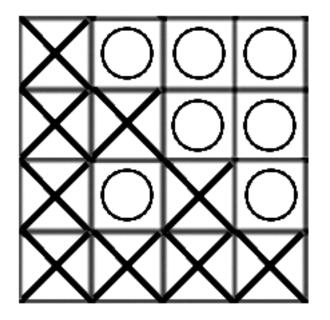
3 2

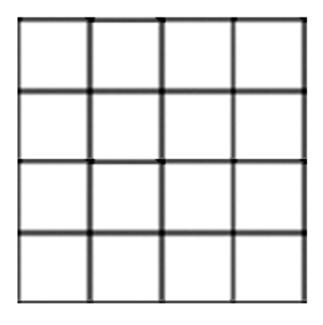


2 3 2

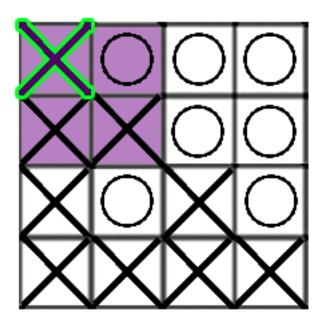
3 3 2

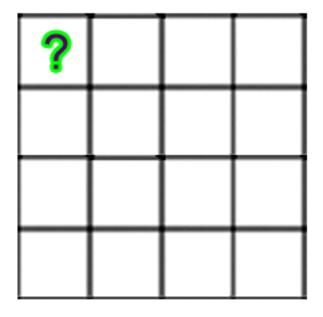
Starting with the grid of X's and O's

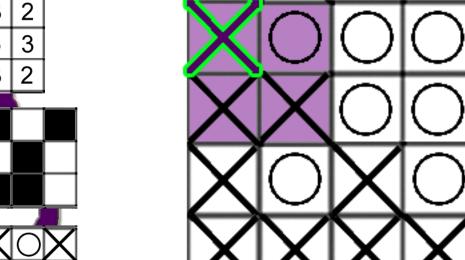


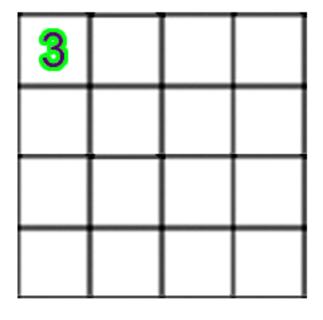


How many X's are in the neighborhood?

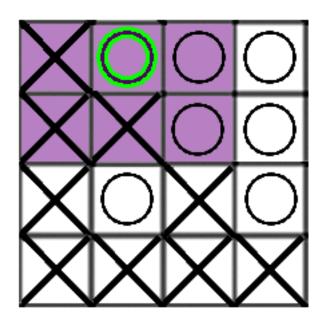


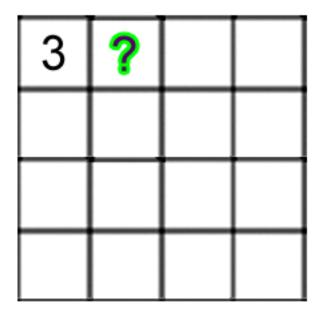


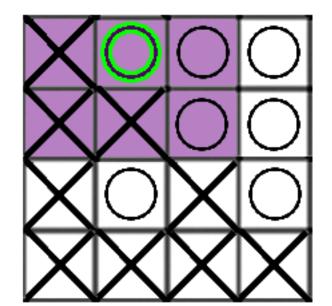


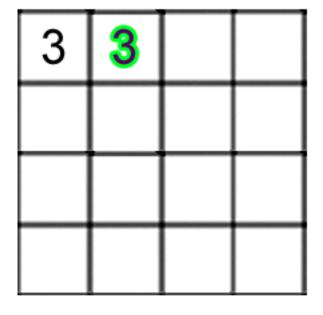


How many X's are in this neighborhood?

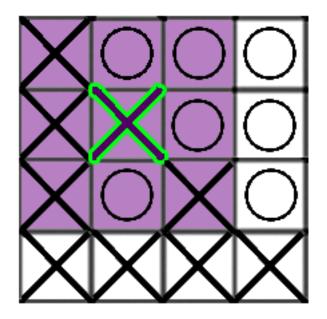




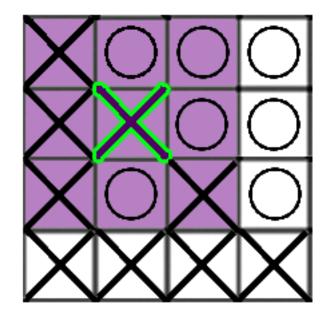




And this neighborhood?

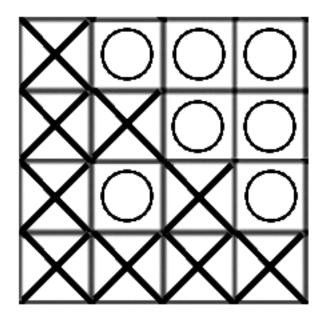


3	3	
	?	



3	3	
	5	

So it's easy to create the puzzle,



3	3	1	0
4	5	2	1
5	7	5	3
3	5	4	3

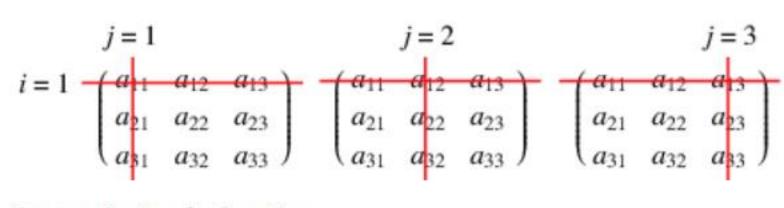
but the fun part is solving it...

Counter example for 2 mod(3)

Given any NxM grid where N=2 mod(3), it is possible to fill the squares with x's and o's such that the clues are all 1's in at least two different ways as follows.



Using Expansion by Minors



For example, for a 3×3 matrix

$$\begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} = a_{11} \begin{vmatrix} a_{22} & a_{23} \\ a_{32} & a_{33} \end{vmatrix} - a_{12} \begin{vmatrix} a_{21} & a_{23} \\ a_{31} & a_{33} \end{vmatrix} + a_{13} \begin{vmatrix} a_{21} & a_{22} \\ a_{31} & a_{32} \end{vmatrix}$$

(from Wolfram MathWorld)